OPTIMUM-STATISTICAL COLLABORATION TOWARDS EFFICIENT BLACK-BOX OPTIMIZATION PURDUE Wenjie Li¹, Chi-Hua Wang¹, and Guang Cheng¹ $\mathbf{V} \mathbf{E} \mathbf{R} \mathbf{S} \mathbf{I} \mathbf{T} \mathbf{Y}_{\mathrm{T}}$ ¹Department of Statistics, Purdue University

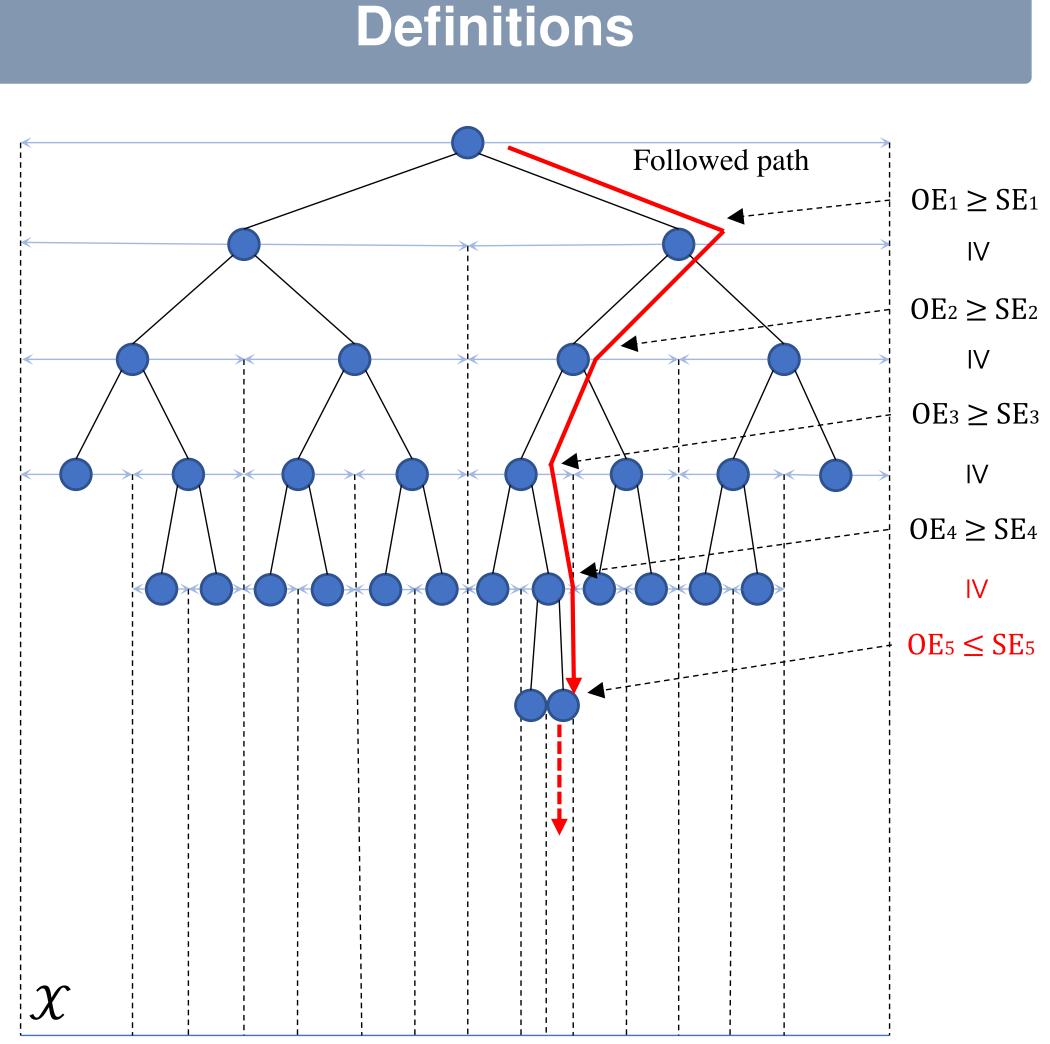
Introduction

Background

- Black-box optimization has gained more and more attention nowadays because of its popularity in application [5, 4]
- Existing works are scattered, and only focus on one specific type of local-smoothness assumption

Contributions

In this paper, we introduce the optimum-statistical collaboration framework with its analysis. In terms of resolution, such an analysis uses a general local smoothness assumption. In terms of uncertainty, the framework inspires us to propose a better quantifier.



Definition 1. (Resolution Descriptor) Define OE_h to be the reso*lution* for each level h, which is a function that upper-bounds the change of f around the optimum and measures the current optimization error, i.e., for any global optimum x^* ,

> $orall h \geq 0, orall x \in \mathcal{P}_{h,i_h^*}, f(x) \geq f(x^*) - \mathtt{OE}_h$, (OE)

where \mathcal{P}_{h,i_h^*} is the node on level h in the partition that contain the global optimum x^* .

Definition 2. (Uncertainty Quantifier) Define $SE_{h,i}(t)$ to be the *uncertainty estimate* for each node $\mathcal{P}_{h,i}$ at time t, which is a function that upper-bounds the current statistical error of the average $\widehat{\mu}_{h,i}(t)$ of the rewards obtained with high probability, i.e., the event

 $\mathcal{A}_t = \left\{ orall (h,i), |\widehat{\mu}_{h,i}(t) - f(x_{h,i})| \leq ext{SE}_{h,i}(t)
ight\}$

is a high probability event.

Algorithm and Results

- A more general local smoothness (resolution descriptor) for different functions and partitions.

certainty, and unspecified algorithm policy.

Algorithm 1 Optimum-statistical Collaboration

- where $\phi(h)$ is a function of the level h, e.g., $\phi(h) = \nu \rho^h$, 1/h

time t, the designed policy to select the optimal node \mathcal{P}_{h_t,i_t} in Algorithm 1 satisfies $f^* - f(x_{h_t,i_t}) \leq a \cdot \max\{ SE_{h_t,i_t}(t), OE_{h_t} \}$, where a > 0is an absolute constant. Then for any $H \in [1, H(n))$ we have the following bound on the expected regret

$$\mathbb{E}[R_n] \leq 2aC\sum_{h=1}^{\overline{H}} (\mathtt{OE}_{h-1})^{-ar{d}} \int_1^{T_h(n)} \max_i \mathtt{SE}_{h,i}(s) ds + \sum_{t=1}^n + a\sum_{\overline{H}+1}^{H(n)} \sum_{i\in\mathcal{I}_h(n)} \int_1^{T_{h,i}(n)} \mathtt{SE}_{h,i}(s) ds + \sqrt{2n\log(4n^3)} + rac{1}{4n^3}$$

where $d := d(a, C, OE_{h-1})$ is the near-optimality dimension defined with respect to a, C, and OE_{h-1} , and $T_h(n) = \max_i T_{h,i}(n)$

Example Regret Bounds for Different $\phi(h)$

Theorem 3: General Regret Bound

- Regret bound when $\phi(h) = \nu \rho^h$
- $\mathbb{E}[R_n] \leq 2\sqrt{2n\log(4n^3)} + C_1 V_{ ext{max}}^{rac{1}{d_1+2}} n^{rac{d_1+1}{d_1+2}} (\log n)^{rac{1}{d_1+2}} + C_2 n^{rac{2d_1+1}{2d_1+4}} \log n^{ra}}$ • Regret bound when $\phi(h) = 2/h$
- $\mathbb{E}[R_n] \leq 2\sqrt{2n\log(4n^3)} + ar{C}_1 V_{ ext{max}}^{rac{1}{2d_2+3}} n^{rac{2d_2+2}{2d_2+3}} (\log n)^{rac{1}{2d_2+3}} + ar{C}_2 n^{rac{2d_2+1}{2d_2+3}} \log n$ SE)

Step 1: Refresh the confidence at some specific time steps to update all $SE_{h,i}(t)$ in the tree. **Step 2:** Find the optimal node \mathcal{P}_{h_t,i_t} at time t that satisfies $OE_{h_t} \leq SE_{h_t,i_t}(t)$ and pull \mathcal{P}_{h_t,i_t} . **Step 3:** If $OE_{h_t} \ge SE_{h_t,i_t}(t)$ after the pull, expand \mathcal{P}_{h_t,i_t} and explore deeper.

 $orall h \geq 0, orall x \in \mathcal{P}_{h,i_h^*}, f(x) \geq f(x^*) - \phi(h)$

$$ext{SE}_{h,i}(t) \equiv c \sqrt{rac{2 \mathbb{V}_{h,i}(t) \log(1/\widetilde{\delta}(t^+))}{T_{h,i}(t)}} + rac{3bc^2 \log(1/\widetilde{\delta}(t^+))}{T_{h,i}(t)}$$
 $ext{Regret Bounds}$

